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DISCRETE MATHEMATICS

COMPUTER SCIENCE & IT

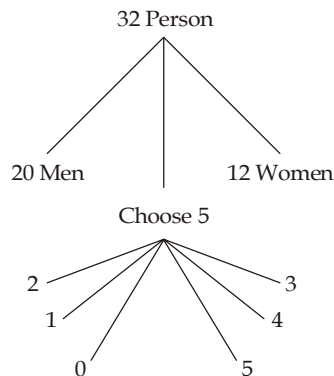
Date of Test : 30/09/2024

ANSWER KEY >

- | | | | | |
|--------|---------|---------|---------|---------|
| 1. (b) | 7. (b) | 13. (a) | 19. (a) | 25. (c) |
| 2. (d) | 8. (a) | 14. (c) | 20. (b) | 26. (a) |
| 3. (b) | 9. (c) | 15. (d) | 21. (c) | 27. (c) |
| 4. (c) | 10. (d) | 16. (b) | 22. (c) | 28. (b) |
| 5. (c) | 11. (a) | 17. (a) | 23. (c) | 29. (c) |
| 6. (c) | 12. (a) | 18. (d) | 24. (c) | 30. (b) |

DETAILED EXPLANATIONS

1. (b)
Given,



We must choose at least 3 women, so, we calculate 3 women, 4 women and 5 women and by addition rule add the results:

$$\begin{aligned}
 &= {}^{12}C_3 \times {}^{20}C_2 + {}^{12}C_4 \times {}^{20}C_1 + {}^{12}C_5 \times {}^{20}C_0 \\
 &= 220 \times 190 + 495 \times 20 + 792 \times 1 = 52492
 \end{aligned}$$

2. (d)
 1. is valid by constructive dilemma.
 2. is valid by destructive dilemma.
 3. is valid by hypothetical syllogism.
 All of the above are known rules of inference.

3. (b)

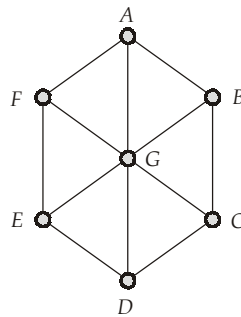
$$\begin{aligned}
 f(x) &= x + 2 & \text{Let } y &= f(x) \\
 y &= x + 2 & \Rightarrow x &= f^{-1}(y) \\
 \Rightarrow x &= y - 2 & & \\
 \Rightarrow f^{-1}(y) &= y - 2 & \text{or } f^{-1}(x) &= x - 2
 \end{aligned}$$

$$\begin{aligned}
 g(3) &= (1 + (3)^2)^{-1} = (1 + 9)^{-1} = \frac{1}{10} \\
 f^{-1} g(3) &= f^{-1}(g(3)) = g(3) - 2 \\
 \Rightarrow f^{-1} g(3) &= \frac{1}{10} - 2 = -1.9
 \end{aligned}$$

4. (c)

Consider each options:

- (a) Null graph of 6 vertices is 1-chromatic so it is correct.
 (b) It is correct because tree with 2 or more vertices is always bichromatic.
 (c) It is incorrect. Consider a wheel graph of 7 vertices.



The chromatic number of graph is 3.

- Color 1 for G
- Color 2 for A, E, C
- Color 3 for F, B, D
- A wheel graph is 3-chromatic when n -vertices are odd and 4-chromatic when n -vertices is even.
- So here $n = 7, \left(\left\lfloor \frac{n}{2} \right\rfloor + 1 \right) = \left(\left\lfloor \frac{7}{2} \right\rfloor + 1 \right) = 4$ which is incorrect because only 3 colors are required to color the above wheel graph.

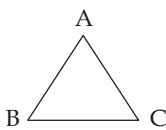
(d) This statement is correct because graph without odd length cycle having atleast 1 edge is bichromatic.

All other statements are true except option (c).

5. (c)

Considering each statements:

- If a graph is bipartite, then its two colourable. Because a bipartite graph can be represented as two groups of vertices such that vertices in same graph are not adjacent. Similarly, statement 2 is equivalent to statement 1.
- If a bipartite graph has a cycle, then it has to be of even length. Graph G is bipartite iff no odd length cycle.



- This graph has a Hamiltonian circuit, but the cycle is of odd length and not bipartite.

$\therefore 3 \not\cong 4$



This graph is bipartite and 2 colorable but does not have Hamiltonian circuit.

So 1, 2 and 4 are equivalent statements.

6. (c)

$$a_n = -5a_{n-1} + 6a_{n-2}$$

$$a_n + 5a_{n-1} - 6a_{n-2} = 0$$

$$x^2 + 5x - 6 = 0$$

$$\begin{aligned}
 x^2 + 6x - x - 6 &= 0 \\
 x(x + 6) - 1(x + 6) &= 0 \\
 (x + 6)(x - 1) &= 0
 \end{aligned}$$

$$x = -6, \quad x = 1$$

$$\begin{aligned}
 \therefore a_n &= A(-6)^n + B \cdot (1)^n \\
 a_n &= A(-6)^n + B
 \end{aligned}$$

7. (b)

$$\text{Number of ways all vowels are together} = \frac{5! \times 4!}{2! 2!} = 720$$

Number of ways not all vowels are together = Total number of permutation - Number of ways all vowels are together

$$= \frac{8!}{2! 2!} - 720$$

$$\Rightarrow \text{Number of not all vowels together} = (10080 - 720) = 9360$$

8. (a)

Hamiltonian cycle for the above graph G is *abcdefa*.

Condition: Each node should be visited exactly once.

9. (c)

The theorem is every finite lattice is bounded but a bounded lattice may not be finite.

10. (d)

Complementary lattice may not have unique complement for every element.

11. (a)

$A \cup B \subseteq A \cap B$ holds true when $A = B$. It is true for empty as well as nonempty sets.

$\Rightarrow |A| = |B|$ is true $|A| \geq 0$ eg. $A = B = \{a, b\}$

Hence $A = \{ \}, B = \{ \}$ "always" is false.

12. (a)

R is reflexive: Since $(a, b) R (a, b)$ for all elements (a, b) because $a = a$ and $b = b$ are always true.

R is symmetric: Since $(a, b) R (c, d)$ and $a = c$ or $b = d$ which can be written as $c = a$ or $d = b$.

So, $(a, b) R (a, b)$ is true.

R is not antisymmetric: Since $(1, 2) R (1, 3)$ and $1 = 1$ or $2 = 3$ true b/c $1 = 1$.

So $(1, 3) R (1, 2)$ but here $2 \neq 3$ so $(1, 2) \neq (1, 3)$.

So, only statement 1 and 2 are correct.

13. (a)

Total number of subset of 5 element = ${}^{25}C_5$

$$= \frac{25 \times 24 \times 23 \times 22 \times 21}{5 \times 4 \times 3 \times 2 \times 1} = 23 \times 22 \times 21 \times 5 = 53130$$

T be a 5 element subset contain no odd number = ${}^{12}C_5$

$$= \frac{12 \times 11 \times 10 \times 9 \times 8}{5 \times 4 \times 3 \times 2 \times 1} = 792$$

So number of 5 element subset with atleast 1 odd number

$$\begin{aligned} T \subseteq S &= {}^{25}C_5 - {}^{12}C_5 \\ &= 53130 - 792 = 52338 \end{aligned}$$

14. (c)

A number is relatively prime to 15 iff it is not divisible by 3 and not divisible by 5.

$$\text{Set of integer from 1 to 1000 divisible by 3} = \left\lfloor \frac{1000}{3} \right\rfloor = 333.$$

$$\text{Set of integer from 1 to 1000 divisible by 5} = \left\lfloor \frac{1000}{5} \right\rfloor = 200.$$

So, number of integer not relatively prime to 15 are

$$\begin{aligned} |A \cup B| &= |A| + |B| - |A \cap B| \\ &= \left\lfloor \frac{1000}{3} \right\rfloor + \left\lfloor \frac{1000}{5} \right\rfloor - \left\lfloor \frac{1000}{15} \right\rfloor \\ &= 333 + 200 - 66 = 467 \end{aligned}$$

So, number of integer relatively prime to 15 are

$$|\overline{A \cup B}| = 1000 - 467 = 533$$

15. (d)

“Not every satisfiable logic is Valid”

$$= \text{Not (every satisfiable logic is Valid)}$$

$$= \text{Not } (\forall x (\text{satisfiable } (x) \Rightarrow \text{Valid } (x))) \quad \text{option (a)}$$

$$= \text{Not } (\forall x (\neg \text{satisfiable} \vee \text{Valid } (x))) \quad \text{option (c)}$$

$$= \exists x (\text{satisfiable } (x) \wedge \neg \text{Valid } (x)) \quad \text{option (b)}$$

Statement (d) says every satisfiable logic is invalid.

So option (d) is not represent given statement.

16. (b)

- I is not D_{42} because the divisor 7 is missing. So, there is no way for I to be isomorphic to $(P\{a,b,c\}, \subseteq)$ as it needs to have 8 divisors but right now it has only 7.
- II is D_{66} a well known boolean algebra and has 8 vertices and its masses diagram will be isomorphic $(P(\{a,b,c\}), \subseteq)$.
- III is not isomorphic even though it looks like D_{70} , it is on the relation \leq , resulting in a chain, which won't be boolean algebra.

17. (a)

$$T(n) - 9T(n-1) + 20T(n-2) = 0$$

$$\text{Let } a_n = T(n)$$

$$\Rightarrow a_n - 9a_{n-1} + 20a_{n-2} = 0$$

$$t^2 - 9t + 20 = 0$$

$$\begin{aligned}
 t^2 + 5t - 4t + 20 &= 0 \\
 t(t - 5) - 4(t - 5) &= 0 \\
 (t - 4)(t - 5) &= 0 \\
 t &= 4, 5
 \end{aligned}$$

Homogenous equation become

$$a_n = c_1 \cdot 5^n + c_2 \cdot 4^n \quad \dots(1)$$

Put $n = 0$ in equation (1)

$$\begin{aligned}
 a_0 &= c_1 \cdot 5^0 + c_2 \cdot 4^0 \\
 -3 &= c_1 + c_2 \quad \dots(2)
 \end{aligned}$$

Put $n = 1$ in equation (1)

$$\begin{aligned}
 a_1 &= c_1 \cdot 5^1 + c_2 \cdot 4^1 \\
 -10 &= 5c_1 + 4c_2 \quad \dots(3)
 \end{aligned}$$

Solving equation (2) and (3) and get c_1 and c_2

$$\begin{aligned}
 (c_1 + c_2 = -3) \times 5 \\
 5c_1 + 4c_2 &= -10
 \end{aligned}$$

$$5c_1 + 5c_2 = -15$$

$$\underline{5c_1 + 4c_2 = -10}$$

$$c_2 = -5 \quad \text{and} \quad c_1 = 2$$

Put value of c_1 and c_2 in eq. (1)

$$a_n = 2 \cdot 5^n - 5 \cdot 4^n$$

18. (d)

- S_1 is correct because connected graph has a Euler circuit if and only if it has number of odd degree vertices is 0.
- A connected graph has a Euler path if and only if it has number of odd degree vertices is either 0 or 2. Therefore a connected graph has Euler path but not euler circuit if and only it has exactly 2 vertices of odd degree therefore S_2 is correct.
- A complete graph of n -vertices contains $n - 1$ degree at each vertex which is greater than $\frac{n}{2}$ for all $n \geq 3$ therefore complete graph has a Hamiltonian circuit. So S_3 is correct
- C_6 is bipartite because any cycle graph with even number of vertices is bipartite. A complete graph with 4 vertices contains complete graph of 3 vertices which contains odd length cycle hence it is not bipartite. So statement S_4 is incorrect.

19. (a)

Maximum and minimum number of component given by:

$$n - K \leq e \leq \frac{(n - K + 1)(n - K)}{2}$$

1. $n - K \leq e$

$$n - e \leq K$$

$$10 - 6 \leq K$$

(\because Minimum number of component)

2. $e \leq \frac{(n - K + 1)(n - K)}{2}$

$$6 \leq \frac{(10 - K + 1)(10 - K)}{2}$$

$$\begin{aligned}
 2 \times 6 &\leq (11 - K)(10 - K) \\
 12 &\leq (10 - K)(11 - K) \\
 12 &\leq K^2 + 110 - 21K \\
 0 &\leq K^2 + 98 - 21K \\
 K^2 + 98 - 21K &= 0 \\
 K &= 14, 7
 \end{aligned}$$

Maximum value of K is 7 because number of components never be larger than nodes.

20. (b)

Since bit are '0' and '1' form. The hamming distance relation on bit has a digraph which will be always an 5-cube where 5 is the number of bits.

- Chromatic number of n -cube = 2 (Since n -cube is always bipartite)

So chromatic number of 5-cube = 2

i.e., '0' = One color
 '1' = Second color

- Diameter of n -cube = n

Diameter of 5 cube = 5

i.e., maximum length between any two vertex.

So ratio $\frac{2}{5} = \frac{X}{Y}$
 $Y - X = 5 - 2 = 3$

21. (c)

Total number of terms = $6 + 1 = 7$

So middle term is 4th term.

$(x + y)^n$ has $(r + 1)$ th term as ${}^n C_r x^{n-r} y^r$.

[(3 + 1)th term] 4th term is

$$\begin{aligned}
 &= {}^6 C_3 \left(\frac{\sqrt{x}}{3} \right)^{6-3} \left(\frac{-3}{x\sqrt{y}} \right)^3 \\
 &= {}^6 C_3 \cdot \left(\frac{(\sqrt{x})^3}{27} \right) \cdot \left(\frac{-27}{x^3 \cdot (\sqrt{y})^3} \right) \\
 &= 20 \cdot \left(\frac{x^{3/2}}{27} \right) \cdot \left(\frac{-27}{x^3 \cdot (y)^{3/2}} \right) \\
 &= -20 \left(\frac{x}{y} \right)^{\frac{3}{2}} \cdot \frac{1}{x^3}
 \end{aligned}$$

22. (c)

I and IV are true. Let's see why IV is true first. We will treat set theory as boolean algebra here, and will demonstrate how to apply this approach.

Given, $S \subseteq R$, which in same as, $\Rightarrow S - R = \emptyset$

The same can be written in boolean algebra as, $\Rightarrow S \wedge R' = 0$

Since we know that \wedge is commutative, $\Rightarrow R' \wedge S = 0$

Now $R' \wedge S = 0$ is same as $R' \wedge (S')' = 0 \Rightarrow R' - S' = \phi$

$\Rightarrow R' \subseteq S'$

Therefore IV is correct.

In order to show that R^{-1} is a subset of S^{-1} , we just need to show that every element in R^{-1} belongs to S^{-1} . So let's assume that R is a subset of S . So if (a, b) is an element of R , then (a, b) belongs to S as well. As (a, b) belongs to R , (b, a) belongs to R^{-1} . Also, (a, b) belongs to S , (b, a) will belong to S^{-1} . Since we can show this presence of every element in R^{-1} in S^{-1} , we see that R^{-1} is a subset of S^{-1} . However II is clearly not true.

23. (c)

This problem corresponds to the number of non-negative integral solution to

$$x_1 + x_2 + x_3 = 10 \text{ with the conditions}$$

$$0 \leq x_1 \leq 10$$

$$0 \leq x_2 \leq 5$$

$$0 \leq x_3 \leq 3$$

Generating functions are required, since the variables have an upper constraint.

Generating function is

$$(1 + x + x^2 + \dots + x^{10}) (1 + x + x^2 + \dots + x^5) (1 + x + \dots + x^3)$$

$$= \left(\frac{1-x^{11}}{1-x} \right) \left(\frac{1-x^6}{1-x} \right) \left(\frac{1-x^4}{1-x} \right)$$

$$= \frac{(1-x^{11})(1-x^6)(1-x^4)}{(1-x)^3}$$

$$= (1-x^4 - x^6 + x^{10}) \sum_{r=0}^{\infty} {}^{3-1+r}C_r \cdot x^r$$

$$= (1-x^4 - x^6 + x^{10}) \sum_{r=0}^{\infty} {}^{r+2}C_r \cdot x^r$$

The coefficient of x^{10} in above generating function is

$${}^{12}C_{10} - {}^8C_6 - {}^6C_4 - {}^2C_0 = 24$$

24. (c)

R^1 is nothing but R itself.

Now, $R^2 = R \circ R$ i.e. composite of R with R

If $(a, b) \in R$, then $(a, c) \in R^2$ iff $(b, c) \in R$.

$$\text{So, } R^2 = \{(1, 1), (2, 1), (3, 1), (4, 1)\}$$

$$R^3 = \{(1, 1), (2, 1), (3, 1), (4, 1)\}$$

$$\begin{aligned} \text{So, } P &= R^1 \cup R^2 \cup R^3 \\ &= \{(1, 1), (2, 1), (3, 1), (4, 1), (4, 2)\} \end{aligned}$$

$$\therefore \text{Cardinality} = 5$$

25. (c)
 “ f is one-one and onto”.
 Negation of this statement will be
 “ f is not one-one or not onto”.
 Now, according to statement R .

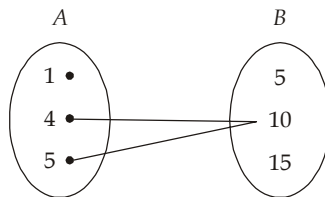
Let,

$$a_1 = 4 \in A$$

$$a_2 = 5 \in A$$

$$f(a_1) = f(a_2) = 10$$

So,

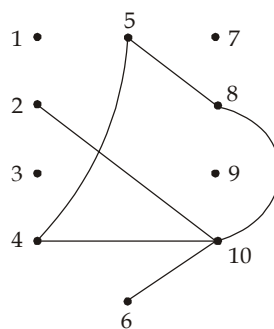


Clearly this is condition of not one-one.
 So, R is correct.
 Now, Q is definition of onto so we have to take negation of this.
 Therefore option (c) is correct answer.

26. (a)
 $(Z, +)$ is a group and $Z \subseteq Q$.
 $(A, +)$ is not a group. Hence it is not a subgroup of $(Q, +)$.
 $(B, +)$ is not a group. Hence it is not a subgroup of $(Q, +)$.

27. (c)
 G is a planar graph. Every planar graph is 4 colorable. Every face is bordered by 3 edges.
 So graph has possibilities of 3 or 4 colors.
 k_3 colored with 3 and k_4 colored with 4 colors.

28. (b)



$$E = \{\{2, 10\}, \{4, 10\}, \{6, 10\}, \{8, 10\}, \{4, 5\}, \{5, 8\}\}$$

\Rightarrow 6 edges are present in G .

29. (c)

$$4.r = 2.e \quad [\text{Planar graph}] \quad \dots(i)$$

$$3.n = 2.e \quad [\text{Cubic graph}] \quad \dots(ii)$$

$$n - e + r = 2 \quad \dots(iii)$$

Substitute (i) and (ii) in (iii)

$$\Rightarrow \frac{2e}{3} - e + \frac{2e}{4} = 2$$

$$\Rightarrow 8e - 12e + 6e = 24$$

$$\Rightarrow 14e - 12e = 24$$

$$\Rightarrow 2e = 24$$

$$\Rightarrow e = 12$$

30. (b)

'd' precedes the set and $\{a, b, h\}$ succeeds the set.

$$\therefore \text{Lower bound} = d$$

$$\text{Upper bound} = \{a, b, h\}$$

