

# POSTAL Book Package

# 2021

## Computer Science & IT

### Objective Practice Sets

#### Programming and Data Structure

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**Q.1** Which of the following C expressions access the  $(i, j)^{\text{th}}$  entry of an  $(m \times n)$  matrix stored in column major order?

- (a)  $n \times (i - 1) + j$  (b)  $m \times (j - 1) + i$   
 (c)  $m \times (n - j) + j$  (d)  $n \times (m - i) + j$

**Q.2** Consider 3 dimensional Array  $A[90][30][40]$  stored in linear array in column major order. If the base address starts at 10, what is the location of  $A[20][20][30]$ ? Assume the first element is stored at  $A[1][1][1]$ .

**Q.3** Consider a 3-heap tree which is similar to 2-heap tree. Every node in 3-heap contains maximum of 3-children. If Array is used to store the element of 3-heap, find the children of node  $i$ ? Assume the first element of array is at 1.

- (a)  $3i, 3i + 1, 3i + 2$  (b)  $3i - 1, 3i, 3i + 1$   
 (c)  $3i + 1, 3i + 2, 3$  (d) None of these

**Q.4** In a compact single dimensional array representation for lower triangular matrices of size  $n \times n$ , non-zero elements of each row are stored one after another, starting from the first row, the index of the  $(i, j)^{\text{th}}$  element of the lower triangular matrix in this new representation is

- (a)  $i + j$  (b)  $i + j - 1$   
 (c)  $j + \frac{i(i-1)}{2}$  (d)  $i + \frac{j(j-1)}{2}$

**Q.5** Consider the following function:

```
int search (int A[ ], int k, int l, int h)
{
    int m;
    if (l == h)
    if (k == A[l]) return l;
        else return -1;
    m =  $\lfloor (l + h)/2 \rfloor$ ;
    if (k  $\leq$  A[m])
```

```
    return search (A, k, l, m);
    else
    return search (A, k, m + l, h);
}
```

Above function is implemented to search a key in the sorted array with binary search concept. Find the index of key 15 returned by the above function, if array has the following elements and  $l = 0, h = 8$  are passed to the function along with array and key.

A	12	14	15	15	15	18	110	120	125
	0	1	2	3	4	5	6	7	8

- (a) 2 (b) 3  
 (c) 4 (d) None of these

**Q.6** Consider a two-dimensional array with elements stored in the form of lower triangular matrix. How many elements must be crossed to read  $A[4, 2]$  from the array  $A[-6, \dots, +8, -6, \dots, +8]$  whose base address is 1000? (Assume elements are stored in row major order).

**Q.7** Consider the following C code

```
int *P, A[3] = {0, 1, 2};
P = A;
*(P + 2) = 5;
P = A++;
*P = 7;
```

What are the values stored in the array A from index 0 to index 2 after execution of the above code?

- (a) 7, 5, 2 (b) 7, 1, 5  
 (c) 0, 7, 5 (d) None of these

**Q.8** Let's look about the algorithm:

```
int temp, j, i;
for (i = 1; i < n; i++)
{
    temp = A[i];
```

```

for (j = i - 1; j ≥ 0 && (A[j] > temp); j--)
    A[j + 1] = A[j];
    A[j] = temp;

```

If the array is in reverse sorted order then time complexities will be

- (a)  $O(n)$  (b)  $O(n \log_2 n)$   
(c)  $O(n^2)$  (d)  $O(\log_2 n)$

**Q.9** Suppose that we have an array of  $n$  data records to sort and that the key of each record has the value 0 or 1. An algorithm for sorting such a set of records require \_\_\_\_\_ running time.

- (a)  $O(1)$  (b)  $O(n)$   
(c)  $O(n^2)$  (d) None of these

**Q.10** Consider an array  $A$  has  $n$ -elements in which every element is less than  $2n$ . What is the running time to check whether the given array has distinct elements?

- (a)  $O(1)$  (b)  $O(n)$   
(c)  $O(n \log n)$  (d)  $O(n^2)$

**Q.11** Given an array with both +ve and -ve numbers. Find the two elements such that their sum is closest to zero

**Ex.:** 60 -10 70 -80 85 gives -80 85

What is the tightest upper bound to solve this problem?

- (a)  $O(n \log n)$  (b)  $O(n^2)$   
(c)  $O(n^3)$  (d)  $O(n)$

**Q.12** What is the output of the following C code? Assume that the address of  $X$  is 2000 (in decimal) and an integer requires four bytes of memory.

```

int main ()
{
    unsigned int x[4][3] = {{1, 2, 3}, {4, 5, 6}, {7, 8, 9}, {10, 11, 12}};
    printf ("%u,%u,%u", X+3, *(X + 2) + 3);
}

```

- (a) 2036, 2036, 2036  
(b) 2012, 4, 2204  
(c) 2036, 10, 10  
(d) 2012, 4, 6

**Q.13** Consider the following C function:

```

#include <stdio.h>
int main(void)
{
    char c[ ] = "ICRBCSIT17";
    char *p = c;
    printf("%s", c+2[p] - 6[p] - 1);
    return 0;
}

```

The output of the program is

- (a) SI (b) IT  
(c) T1 (d) 17

**Q.14** The output of the following program is

```

main ( )
{
    static int x[ ] = {1, 2, 3, 4, 5, 6, 7, 8};
    int i ;
    for (i = 2; i < 6; ++i)
        x[x[i]] = x[i];
    for (i = 0; i < 8; ++i)
        printf("%d", x[i]);
}

```

- (a) 1 2 3 3 5 5 7 8 (b) 1 2 3 4 5 6 7 8  
(c) 8 7 6 5 4 3 2 1 (d) 1 2 3 5 4 6 7 8

**Q.15** Which of the following is true?

- (a) In sorted array of ' $n$ ' distinct elements, deletion of an element from beginning takes  $O(\log n)$  time.  
(b) In sorted array of ' $n$ ' distinct elements, insertion of an element takes  $O(\log n)$  time.  
(c) In sorted array of ' $n$ ' distinct elements, finding  $i^{\text{th}}$  largest element take  $O(1)$  time.  
(d) In unsorted array of ' $n$ ' distinct elements, insertion of an element take  $O(\log n)$  time.

**Q.16** Consider the function given below, which should return the index of first zero in input array of length ' $n$ ' if present else return -1.

```

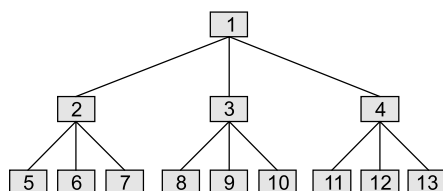
int index of zero (int[ ] array, int n) {
    for (int i = 0; i < n; i++)
        if (array[i] == 0)
            return i;
    return -1;
}

```

For column major order

$$\begin{aligned} \text{loc}(A(i, j, k)) &= \text{Base Address} + (i-1)r_2r_3 + (K-1)r_2 + (j-1) \\ &= 10 + 19 * (30)(40) + 29 * (30) + 19 = 23699 \end{aligned}$$

3. (b)



Children of 4:

$$4 * 3 - 1 = 11$$

$$4 * 3 = 12$$

$$4 * 3 + 1 = 13$$

$\therefore (i * 3) - 1, i * 3$  and  $(i * 3) + 1$  are children of  $i$ .

4. (c)

The number of elements to be skipped to reach

$$\text{to } i^{\text{th}} \text{ row} = \frac{i(i-1)}{2} \text{ to reach to } j^{\text{th}} \text{ column} =$$

$$\frac{i(i-1)}{2} + j.$$

5. (a)

$$\begin{array}{ccccccc} (\ell) & & & (m) & & & (h) \\ 12 & 14 & 15 & 15 & 15 & 18 & 110 & 120 & 125 \end{array}$$

$$\begin{array}{ccccccc} (\ell) & & (m) & & (h) \\ 12 & 14 & 15 & 15 & 15 \end{array}$$

$$\begin{array}{ccc} (\ell) & (m) & (h) \\ 12 & 14 & 15 \end{array}$$

$$15 \Rightarrow l == h \text{ is true and } k == A[l].$$

$$[l = h = 2]$$

$\therefore$  Option (a) is correct.

6. (63)

The given lower triangular matrix can be represented as

$$\begin{array}{ccccccc} & -6 & -5 & -4 & \dots & \dots & +8 \\ -6 & a_{11} & & & & & \\ -5 & a_{21} & a_{22} & & & & \\ -4 & a_{31} & a_{32} & a_{33} & & & \\ \cdot & \cdot & \cdot & & & & \\ \cdot & \cdot & \cdot & & & & \\ \cdot & \cdot & \cdot & & & & \\ \cdot & \cdot & \cdot & & & & \\ +8 & a_{81} & a_{82} & \dots & \dots & \dots & a_{88} \end{array}$$

Let  $(i, j)$  be the element to be accessed.

We must cross upto  $(i-1)^{\text{th}}$  row.

Number of elements upto  $(i-1)^{\text{th}}$  row or  $10^{\text{th}}$  row

$$= 1 + 2 + 3 + \dots + [(i-1) - (l_{bi}) + 1]$$

$[l_{bi} \rightarrow \text{lower bound of } i]$

$$= 1 + 2 + 3 + \dots (3 - (-6) + 1)$$

$$= 1 + 2 + 3 + \dots + (10)$$

$$= \frac{10 \times 11}{2} = 55$$

In  $i^{\text{th}}$  row we must cross  $(j - l_{bj})$  elements.

$[l_{bj} \rightarrow \text{lower bound of } j]$

$$= 2 - (-6) = 8$$

$\therefore$  In total  $= 55 + 8 = 63$  elements need to be crossed.

7. (d)

$P = A++$ ; produces compiler error.

So execution of the given code is not possible.

$A++$  asks the compiler to change the base address of an array, but compiler knows  $A$  is array hence once it is declared, compiler will not allow to change the address.

8. (c)

In this programme first for loop will run  $n$ -times and second for loop also run  $n$ -times, because our array is reverse sorted then second loop will also run and the total time complexity for reverse sorted order will be  $O(n^2)$ .

9. (b)

Using counting sort, it takes linear time.

10. (b)

Using counting sort, single scan will identify if there exist any repeated element in the given array. Therefore, it takes  $O(n)$  time.

11. (a)

1. First sort elements  $\rightarrow n \log n$ .

2. Add  $(i)$  and in temp at last and before that set +ve closest = max and -ve closest-min  
 $\text{temp} = A[i] + A[j]$   
 if  $(\text{temp} > 0)$   
 {  
   if  $(\text{temp} < \text{positive closest})$   
   positive closest = temp;